1. Introduction
Congestion avoidance mechanisms can be broadly divided into detection and control. Congestion can be detected by monitoring the queuing delays and packet drops at the nodes. Congestion control is generally achieved by regulating the packet-sending rate at the source node of the connection. The source node infers the congestion information along the connection route for effective congestion control. To enable the source node to control packet-sending rate, the intermediate routers proactively send the observed congestion levels to the participating sources. Another alternative is a passive mechanism where the source node infers the congestion level by measuring parameters like end-to-end packet delays. A hybrid of active and passive approach is used in the Internet. In a wired network, the packet drop information accurately conveys the congestion information since packet drops are primarily due to overflow of queues at intermediate nodes. Packet drops due to channel error or ineffective MAC protocol is very rare. The end-to-end packet delays are primarily due to the intermediate queuing delays since the transmission time is low.

For example, packet drops can be caused due to a high channel error rate. The end-to-end delays can be skewed due to the wireless MAC layer transmission scheme, which can significantly increase the transmission time of a packet.

2. Nature of Cross Layer Design
TCP is a reliable connection-oriented byte-stream protocol which performs congestion control dynamically during the data communication process. TCP congestion control is performed on an end-to-end basis.

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ABSTRACT
Multi-hop Wireless Networks (MHWNs) are anticipated to play an important role at the edge of the Internet, enabling a large number of innovative applications. The great demand for capacity from a large number of users and applications, coupled with the sparse bandwidth available on the wireless channel, place particular emphasis on effective congestion management approaches. Effective and well-studied algorithms for congestion control at the transport layer exist in wired networks. However, for a number of reasons, these approaches do not translate directly to wireless environments. In this chapter, we first describe the problem of congestion control in MHWNs, and discuss approaches for solving it. The presentation is organized into two components: (1) a review of the causes of congestion and algorithms for congestion avoidance in MHWNs at different layers of protocol stack; and (2) a review of analytical models for the rate control problem and their use for congestion control.
The receiver provides an acknowledgement (ACK) feedback back to the sender. Relying on the information provided by ACKs, the sender can detect which packets are lost during transmission over the communication path.

TCP uses additive increase and multiplicative decrease strategy for its window adjustment according to network conditions. The connection is initiated with window size equal to one packet (1 MSS—Maximum Segment Size). Then, cwnd is increased exponentially for every non-duplicate ACK reception until the Slow Start Threshold (ssthresh) is reached. Prior the connection establishment, ssthresh is set to an initial value, which depends on the implementation of the protocol stack, and then adjusted on the basis of the estimate of the network capacity. This technique is called slow start phase.

When the threshold is reached, TCP enters congestion avoidance phase. The window is increased linearly by one packet for each received ACK. The window growth in this phase is limited to a maximum window size, negotiated between sender and receiver during connection establishment and then updated on the fly during the communication process (the receivers advertised window). There are two ways for TCP sender to detect data loss occurred on the communication link: reception of duplicate ACKs (dupacks) and timeout occurrence. In the first case, a dupack is generated by the receiver upon reception of an out-of-order.

3. Research challenges
3.1.1 Band Sensing

One of the primary requirements of Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks is their ability to scan the band and identify vacant channels available for opportunistic transmission. As the primary user network is physically separate from the secondary user network, the secondary users do not get any direct feedback from primary users regarding their transmission. The secondary users have to depend on their own individual or cooperative sensing ability to detect primary user transmissions. Since the primary users can be spread across a huge geographical area, sensing the entire spectral band accurately is a challenging task. The secondary users have to rely on weak primary transmission signals to estimate their presence.
a decentralized network, where Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks Secondary user Primary user Secondary base station Primary base station.

### 3.3.1. Attacks on Congestion Control Networks

Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks any activity that results in (a) unacceptable interference to the licensed primary users or (b) missed opportunities for secondary users. An attack is considered strong if it involves a minimal number of adversaries performing minimal operations but causing maximum damage/loss to the primary and or secondary users in the network. We describe attacks on five layers in the protocol stack, namely, the physical layer, link layer, network layer, transport layer and application layer.

#### 3.3.2. Physical Layer Attacks

The physical layer is the most basic layer in the protocol stack. It provides the means of transmitting raw signals over the transmission medium. The physical layer determines the bit rate, channel capacity, bandwidth and maximum throughput of the connection. In Congestion Control networks, the physical layer has the capability to transmit at various frequencies across most of the Band. Therefore, when transmission from one frequency band is switched to another frequency band, the switching process attacks on Congestion Control Networks incurs considerable delay in the physical layer of Congestion Control networks.

#### 3.3.3. Intentional Jamming Attack

This is one of the most basic types of attack that can be performed by secondary users.

![Fig. 6. International Jamming of connecting nodes](image)

The attack can be further amplified by using high transmit power, transmitted in several bands. Although simple energy-based detection and triangulation techniques can be used to detect this attack, the time it takes to pinpoint and ban the malicious user impacts severely on the network performance.

#### 3.3.4. Primary Receiver Jamming Attack

The attack is caused when a malicious entity closer to the victim primary receiver participates in a collaborative protocol and requests transmissions from other secondary users to be directed towards the malicious user.

#### 3.3.5. Sensitivity Amplifying Attack

This leads to frequent false detections and missed opportunities for the secondary users. Operating in multiple bands to incur missed opportunities and render Band usage inefficient.

### 3.3.6. Overlapping Secondary User Attack

In both centralized and distributed Congestion Control networks, multiple secondary networks may coexist over the same geographical region.

In such cases, transmissions from malicious entities in one network can cause harm to the primary and secondary users of the other network. This type of attack is hard to prevent because the malicious entities may not be under the direct control of the secondary base station/users of the victim network.

#### 3.3.7. Link Layer Attacks

The link layer provides Security Issues in Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Network the functional means to allow fragmentation of data, error correction and modulation. The medium access control (MAC) layer is one of the important sub layers of the link layer, which controls channel assignment.

#### 3.3.8. Asynchronous Sensing Attack

Instead of synchronizing the sensing activity with other secondary users in the network, a malicious secondary user may transmit asynchronously when other secondary users are performing sensing operations. If the base station or other secondary users consider this as a transmission from a primary user, then this could result in missed opportunities. This attack can be made more efficient by transmitting only during sensing periods.

#### 3.3.9. False Feedback Attack

False feedback from one or a group of malicious users could make other secondary users take inappropriate actions and violate the goals of the protocol. Primary base station’s transmission range. A similar attack is possible in centralized Congestion Control radio network.

#### 3.3.11. Network Layer Attacks

While the data link layer is responsible for node to node (one-hop) packet delivery, the network layer is responsible for end-to-end (source-to-destination) packet delivery. The network layer provides functional means for performing routing, flow control and ensuring quality of service (QoS). Some of the routing protocols used in wireless environment are for example dynamic source routing (DSR) and ad-hoc on demand distance vector (AODV) routing.

A malicious node in the path can disrupt routing by either broadcasting incorrect routing information to its neighbors or by redirecting the packets in the wrong direction. Several routing attacks have been discovered in wireless ad-hoc networks, most of the attacks can be classified into two categories: routing disruption attacks and resource consumption attacks. Some of the examples of routing attacks are the black hole attack where the malicious node attracts packets from every other node and drops all the packets, the gray hole attack where the malicious node selectively drops the packets, the worm hole attack where the malicious user uses two pairs of nodes with a private connection between the two pairs.

The worm hole attack is a dangerous attack since it can prevent route discovery where the source and the destination are more than two hops away. Most of these attacks are prevented by using secure on-demand routing protocols like secure AODV, which use cryptographic mechanisms to guarantee integrity of routing information and authenticity of nodes.

### 4.1.1. Future Directions

We provide some future directions that need to be taken to make secure Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks against both accidental and intentional attacks. Most of our Security Issues in Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks proposed solutions are easy to implement (for example, using existing security protocols).
Security services provided in cellular, WLAN and wireless ad-hoc networks can be applied to Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks as well. It is the last hop between the wireless base stations and the wireless terminals that needs to be protected over the air. Confidentiality is provided by using the confidentiality algorithm known as f8 and the secret cipher key (CK) that is exchanged as a part of the AKA process. Integrity is provided by using the integrity algorithm, f9 and the integrity key (IK). A block cipher known as KASUMI is the building block of both f8 and f9 algorithms. KASUMI operates on 64 bit blocks and uses a 128-bit secret key. Decentralized Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks could use security mechanisms employed in ad-hoc wireless networks.

4.1.3. Using Cryptographic Primitives

Primary user identification is very important for both centralized and decentralized Congestion Control networks. We recently proposed a digital signature based primary user identification mechanism that can be used by secondary users to distinguish malicious transmissions from primaries.

4.1.4. Reactive Security Mechanisms

Reactive security mechanisms that detect malicious activity in Cross Layer Design, For example, mechanisms that can detect unusually high Band handoffs is useful to prevent jamming and Band handoff attacks.

4.1.5. Band Aware Approach

There are two ways to handle Band mobility and associated delays. One is to make Band sensing, analyzing and handoff process fast and transparent to the higher layer protocols. However, Band sensing and handoff processes are in their infant stages and it will take a long time for such approaches to materialize. Another approach is a cross-layer methodology to incorporate Band mobility as state information in protocols operating in upper layers. Although this approach increases cross-layer dependencies on transport layer should consider the effect of Band handoff on the round trip time and correspondingly adjust the retransmission window.

4.1.6. Use Light-weight Security Protocols and Primitives

Light-weight security protocols need to be developed for power/resource constrained environments.

5. Conclusions

The main motivation behind Cross Layer design has been to increase Band utilization by allowing the unlicensed (secondary) users to opportunistically access the frequency band actually owned by the licensed (primary) user. In contrast to other network security architectures, the users are categorized into two distinct classes: primary users and secondary users. We also discussed various security aspects such as authentication and authorization of users, confidentiality and integrity of communication as well as identification and non repudiation of Congestion Control user devices. Some reliability issues that are inherent in cognitive networks were examined. Through these attacks we showed that the fundamental idea behind Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks (to have self-aware networks that offer resilient services and keep the intruders out of it simply by cognition) is yet fulfilled. Finally, we suggested some future directions that need to be taken to make the protocols that are employed in Cross Layer Design Approach for avoid Congestion Control in Wireless Mobile AD-HOC Networks.

6. References.

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